

On second-order nonlinearities of some \mathcal{D}_0 type bent functions

Sugata Gangopadhyay
(Joint work with Brajesh Kumar Singh)

Department of Mathematics, Indian Institute of Technology Roorkee
gsugata@gmail.com

Central European Conference on Cryptology
Bedlewo, Poland
June 10–12, 2010

Boolean functions

- ▶ A Boolean function f on n variables is a mapping :

$$f : \{0, 1\}^n \rightarrow \{0, 1\}.$$

- ▶ Boolean function can be represented by a truth table.

$$f : \{0, 1\}^4 \rightarrow \{0, 1\}$$

x_4	x_3	x_2	x_1	f
0	0	0	0	0
0	0	0	1	1
0	0	1	0	1
0	0	1	1	1
0	1	0	0	1
0	1	0	1	1
0	1	1	0	0
0	1	1	1	0
1	0	0	0	0
1	0	0	1	0
1	0	1	0	0
1	0	1	1	0
1	1	0	0	0
1	1	0	1	0
1	1	1	0	0
1	1	1	1	0

- ▶ Also represented by the 2^n -length string $[f(0), f(1), \dots, f(2^n - 1)]$
- ▶ The number of all n -variable Boolean functions is 2^{2^n} .

Boolean functions

- ▶ Let \mathbb{F}_2 be the prime field of characteristic 2.
- ▶ A Boolean function f is a function from \mathbb{F}_2^n to \mathbb{F}_2 .
- ▶ Alternatively Boolean functions can be thought of as functions from \mathbb{F}_2^n to \mathbb{F}_2 .
- ▶ Let \mathcal{B}_n be the set of all Boolean functions on n variables.

Algebraic Normal Form (ANF)

- ▶ The algebraic normal form (ANF) of $f \in \mathcal{B}_n$ is

$$f(x_1, x_2, \dots, x_n) = \sum_{\mathbf{a}=(a_1, \dots, a_n) \in \mathbb{F}_2^n} \mu_{\mathbf{a}} \left(\prod_{i=1}^n x_i^{a_i} \right),$$

where $\mu_{\mathbf{a}} \in \mathbb{F}_2$.

- ▶ The algebraic degree of f ,
 $\deg(f) := \max\{wt(\mathbf{a}) : \mu_{\mathbf{a}} \neq 0, \mathbf{a} \in \mathbb{F}_2^n\}$.

Nonlinearity

- ▶ $dist : \mathcal{B}_n \times \mathcal{B}_n \longrightarrow \mathbb{Z}$ defined by $dist(f, g) = |\{x \in \mathbb{F}_2^n : f(x) \neq g(x)\}|$, for all $f, g \in \mathcal{B}_n$, is said to be the Hamming distance between f and g .
- ▶ Nonlinearity of $f \in \mathcal{B}_n$ is defined as $nl(f) = \min_{l \in \mathcal{A}_n} \{dist(f, l)\}$ where \mathcal{A}_n is the set of affine functions on n variables.
- ▶ Alternatively the nonlinearity of f is its distance from $RM(1, n)$, the Reed-Muller code of order 1 and size 2^n .

Nonlinearity and Walsh Transformation

- ▶ The Walsh transform $f \in \mathcal{B}_n$ at $\lambda \in \mathbb{F}_2^n$ is defined as follows:

$$W_f(\lambda) = \sum_{x \in \mathbb{F}_2^n} (-1)^{f(x) + \lambda \cdot x}.$$



$$nl(f) = 2^{n-1} - \frac{1}{2} \max_{\lambda \in \mathbb{F}_2^n} |W_f(\lambda)|.$$

It is to be noted that the Walsh spectrum of $f \in \mathcal{B}_n$ can be computed in time $O(n2^n)$ and hence the nonlinearity.

Upper bound of nonlinearity

- ▶ Parseval's identity

$$\sum_{\lambda \in \mathbb{F}_2^n} W_f(\lambda)^2 = 2^{2n}$$

- ▶ $|W_f(\lambda)| \geq 2^{n/2}$, which implies $nl(f) \leq 2^{n-1} - 2^{\frac{n}{2}-1}$.

Bent functions: functions with maximum nonlinearity

- ▶ A Boolean function $f \in \mathcal{B}_n$, where n is even is said to be bent if and only if $|W_f(\lambda)| = 2^{n/2}$ for all $\lambda \in \mathbb{F}_2^n$.
- ▶ From this it follows that bent functions have maximum nonlinearity namely $2^{n-1} - 2^{\frac{n}{2}-1}$ for even n .

Nonlinearity to nonlinearity profile

- ▶ Suppose f is a Boolean function on n variables. For every non-negative integer $r \leq n$, we denote by $nl_r(f)$ the r th-order nonlinearity of f , which is the minimum Hamming distance of f and all functions of algebraic degree at most r .
- ▶ Alternatively the r th-order nonlinearity of an n variable Boolean function f is its distance from $RM(r, n)$, the r order Reed-Muller code of size 2^n .
- ▶ The sequence of values $nl_r(f)$, for r ranging from 1 to $n - 1$, is said to be the nonlinearity profile of f .
- ▶ Unlike the first-order nonlinearity there is no fast algorithm to determine second or higher-order nonlinearities.

Lower bounds on nonlinearity profile

- ▶ (Carlet 2008) C. Carlet, Recursive lower bounds on the nonlinearity profile of Boolean functions and their applications, IEEE Trans. Inform. Theory 54 (3) (2008) 1262-1272.
- ▶ (Fouquet and Tavernier 2008) R. Fourquet and C. Tavernier, An improved list decoding algorithm for the second order Reed-Muller codes and its applications, Des. Codes Cryptogr. 49 (2008) 323-340.

Derivatives of Boolean function

- ▶ The derivative of $f \in \mathcal{B}_n$ with respect to $a \in \mathbb{F}_2^n$ is defined by

$$D_a f(x) := f(x) + f(x + a)$$

- ▶ The r th-order derivative of f with respect to V is defined by

$$D_V f(x) := D_{a_1} \dots D_{a_r} f(x)$$

Where V be an r -dimensional subspace of \mathbb{F}_2^n generated by a_1, \dots, a_r .

Proposition 2 (Carlet 2008)

- ▶ Let f be n variable Boolean function and r be a positive integer smaller than n , then we have

$$nl_r(f) \geq \frac{1}{2} \max_{a \in \mathbb{F}_2^n} nl_{r-1}(D_a f)$$

- ▶ In particular, for $r = 2$

$$nl_2(f) \geq \frac{1}{2} \max_{a \in \mathbb{F}_2^n} nl(D_a f).$$

Proposition 3 and Corollary 2 (Carlet 2008)

- ▶ Let f be any n variable Boolean function and r be a positive integer smaller than n , Then we have
$$nl_r(f) \geq 2^{n-1} - \frac{1}{2} \sqrt{2^{2n} - 2 \sum_{a \in \mathbb{F}_2^n} nl_{r-1}(D_a f)}.$$
- ▶ Let f be any n variable function and r a positive integer smaller than n . Assume that, for some nonnegative integers M and m , we have $nl_{r-1}(D_a f) \geq 2^{n-1} - M2^m$ for every nonzero $a \in \mathbb{F}_2^n$. Then

$$\begin{aligned} nl_r(f) &\geq 2^{n-1} - \frac{1}{2} \sqrt{(2^n - 1)M2^{m+1} + 2^n} \\ &\approx 2^{n-1} - \sqrt{M2^{\frac{n+m-1}{2}}} \end{aligned} \tag{1}$$

Quadratic Boolean functions

- ▶ Suppose $g \in \mathcal{B}_n$ is a quadratic function. The bilinear form associated with g is defined by
$$B(x, y) = g(0) + g(x) + g(y) + g(x + y).$$
- ▶ The kernel of $B(x, y)$ is the subspace of \mathbb{F}_2^n defined by

$$\mathcal{E}_g = \{x \in \mathbb{F}_2^n : B(x, y) = 0 \text{ for all } y \in \mathbb{F}_2^n\}.$$

Quadratic Boolean functions

- ▶ Suppose $g \in \mathcal{B}_n$ is a quadratic function. The kernel of $B(x, y)$

$$\mathcal{E}_g = \{a \in \mathbb{F}_2^n : D_a g = \text{constant}\}.$$

- ▶ A. Canteaut, P. Charpin and G. M. Kyureghyan, A new class of monomial bent functions, *Finite Fields and their Applications* 14 (2008) 221-241.

Walsh spectrum of quadratic Boolean functions

If $g : \mathbb{F}_2^n \rightarrow \mathbb{F}_2$ is a quadratic boolean function and $B(x, y)$ is the quadratic form associated to it, then the Walsh Spectrum of g depends only on the dimension, k , of the kernel, \mathcal{E}_g , of $B(x, y)$. The weight distribution of the Walsh spectrum of g is:

$W_g(\mu)$	number of μ
0	$2^n - 2^{n-k}$
$2^{(n+k)/2}$	$2^{n-k-1} + (-1)^{f(0)} 2^{(n-k-2)/2}$
$-2^{(n+k)/2}$	$2^{n-k-1} - (-1)^{f(0)} 2^{(n-k-2)/2}$

- ▶ F. J. MacWilliams and N. J. A. Sloane, The theory of error correcting codes, North-Holland, Amsterdam, 1977.

Lower bounds of second-order nonlinearities of cubic Boolean functions

- ▶ If f is a cubic Boolean function the $D_a f$ is at most quadratic.
- ▶ It possible to get good estimates of the nonlinearities of $D_a f$ for all $a \in \mathbb{F}_2^n$ to obtain estimates of the lower bounds of second-order nonlinearities of cubic Boolean functions.
- ▶ This technique is used in several recent papers for cubic bent functions.

Lower bounds of second-order nonlinearities of cubic Boolean functions

- ▶ If $f(x, y) = \text{Tr}_1^p(xy^{2^i+1})$, where $x, y \in \mathbb{F}_{2^p}$, $n = 2p$, $n \geq 6$ and i is an integer such that $1 \leq i < p$, $\gcd(2^p - 1, 2^i + 1) = 1$ and $\gcd(i, p) = e$, then

$$nl_2(f) \geq 2^{n-1} - \frac{1}{2} \sqrt{2^{2(\frac{3n}{2}+e)} - 2^{2(\frac{3n}{4}+\frac{e}{2})} + 2^n(2^{2(\frac{n}{4}+\frac{e}{2})} - 2^e + 1)}.$$

- ▶ S. Gangopadhyay, S. Sarkar and R. Telang, On the lower bounds of the second order nonlinearities of some Boolean functions, Information Sciences 180 (2010) 266-273.

Lower bounds of second-order nonlinearities of cubic Boolean functions

$n = 2p$	6	10	12
i	1, 2	1, 2, 3, 4	2, 4
$e = \gcd(i, p)$	1	1	2
Lower bounds in (Gangopadhyay et al.)	15	378	1524
Hamming distances in (Fourquet et al.)	18	400	1760

Construction \mathcal{D}_0 type bent functions

- ▶ Let $n = 2p$, π is a permutation on \mathbb{F}_2^p . $f(x, y) = \pi(y) \cdot x$ is a Maiorana-McFarland type bent.
- ▶ Following is the \mathcal{D}_0 type bent constructed by Carlet:

$$h(x, y) = x \cdot \pi(y) + \prod_{j=1}^p (x_j + 1)$$

where $x = (x_1, \dots, x_n)$

- ▶ C. Carlet, Two new classes of bent functions, in Proc. EUROCRYPT '93, LNCS vol. 765, Springer, 1994, pp. 77-101.

Walsh transforms of derivatives of \mathcal{D}_0 type bent functions

Let $h(x, y) = f(x, y) + g(x)$, where $n = 2p$, $x, y \in \mathbb{F}_2^p$,
 $f(x, y) = x \cdot \pi(y)$, $g(x) = \prod_{i=1}^p (x_i + 1)$ and π is a permutation
on \mathbb{F}_2^p then

- ▶ The Walsh transform of $D_{(a,b)}h$ at $(\mu, \eta) \in \mathbb{F}_2^p \times \mathbb{F}_2^p$ is

$$W_{D_{(a,b)}h}(\mu, \eta) = W_{D_{(a,b)}f}(\mu, \eta) - 2[(-1)^{\mu \cdot a} + (-1)^{\eta \cdot b}]W_{a \cdot \pi}(\eta)$$

- ▶ $|W_{D_{(a,b)}h}(\mu, \eta)| \leq |W_{D_{(a,b)}f}(\mu, \eta)| + 4|W_{a \cdot \pi}(\eta)|.$

Proof outline

Let $h(x, y) = f(x, y) + g(x)$, $g(x) = \prod_{i=1}^p (x_i + 1)$ and $(a, b) \in \mathbb{F}_2^p \times \mathbb{F}_2^p$, with $a \neq 0$. Then

$$g(x) + g(x + a) = \begin{cases} 1, & \text{if } (x, y) \in (\{0\} \times \mathbb{F}_2^p) \cup (\{a\} \times \mathbb{F}_2^p), \\ 0, & \text{otherwise.} \end{cases}$$

The Walsh transform of $D_{(a,b)}h$ at $(\mu, \eta) \in \mathbb{F}_2^p \times \mathbb{F}_2^p$ is

$$\begin{aligned} W_{D_{(a,b)}}h(\mu, \eta) &= \sum_{(x,y) \in \mathbb{F}_2^p \times \mathbb{F}_2^p} (-1)^{f(x+a,y+b)+f(x,y)+g(x+a)+g(x)+\mu \cdot x + \eta \cdot y} \\ &= \sum_{(x,y) \in \mathbb{F}_2^p \times \mathbb{F}_2^p} (-1)^{f(x+a,y+b)+f(x,y)+\mu \cdot x + \eta \cdot y} \\ &\quad - 2 \sum_{(x,y) \in \{0,a\} \times \mathbb{F}_2^p} (-1)^{f(x+a,y+b)+f(x,y)+\mu \cdot x + \eta \cdot y} \end{aligned}$$

Proof outline

$$\begin{aligned} &= W_{D(a,b)} f(\mu, \eta) - 2 \left[\sum_{y \in \mathbb{F}_2^p} (-1)^{f(0,y+b)+f(a,y)+\mu \cdot a + \eta \cdot y} \right. \\ &+ \left. \sum_{y \in \mathbb{F}_2^p} (-1)^{f(a,y+b)+f(0,y)+\eta \cdot y} \right] \\ &= W_{D(a,b)} f(\mu, \eta) - 2 \left[(-1)^{\mu \cdot a} \sum_{y \in \mathbb{F}_2^p} (-1)^{a \cdot \pi(y) + \eta \cdot y} \right. \\ &+ \left. (-1)^{\eta \cdot b} \sum_{y \in \mathbb{F}_2^p} (-1)^{a \cdot \pi(y+b) + \eta \cdot (y+b)} \right] \\ &= W_{D(a,b)} f(\mu, \eta) - 2 \left[(-1)^{\mu \cdot a} + (-1)^{\eta \cdot b} \right] W_{a \cdot \pi}(\eta) \end{aligned}$$

Thus, $|W_{D(a,b)} h(\mu, \eta)| \leq |W_{D(a,b)} f(\mu, \eta)| + 4 |W_{a \cdot \pi}(\eta)|$

Main Theorem

Let $h(x, y) = \text{tr}_1^p(xy^{2^i+1}) + \prod_{i=1}^p(x_i + 1)$, where $n = 2p$,
 $x, y \in \mathbb{F}_2^p$, i is integer such that $1 \leq i \leq p$,
 $\gcd(2^i + 1, 2^p - 1) = 1$, and $\gcd(i, p) = e$, then

$$nl_2(h) \geq 2^{2p-1} - \frac{1}{2} \sqrt{2^{3p+e} + 2^{2p}(1 - 2^e) + 5(2^{\frac{5p+e}{2}} - 2^{\frac{3p+e}{2}})}.$$

Proof outline

Let $h(x, y) = \text{tr}_1^p(xy^{2^i+1}) + \prod_{i=1}^p(x_i + 1)$, where $n = 2p$,
 $x, y \in \mathbb{F}_2^p$, i is integer such that $1 \leq i \leq p$,
 $\gcd(2^i + 1, 2^p - 1) = 1$, and $\gcd(i, p) = e$, then nonlinearity of
 $D_{(a,b)}h$ is

$$nl(D_{(a,b)}h) \geq \begin{cases} 2^{2p-1} - 2^{p+e-1}, & \text{if } a = 0 \text{ and } b \neq 0, \\ 2^{2p-1} - 2^{p+e-1} - 2^{\frac{p+e+2}{2}}, & \text{if } a \neq 0 \text{ and } b \neq 0, \\ 2^{2p-1} - 2^{\frac{3p+e-2}{2}} - 2^{\frac{p+e+2}{2}}, & \text{if } a \neq 0 \text{ and } b = 0. \end{cases}$$

Proof outline



$$\begin{aligned} & \sum_{(a,b) \in \mathbb{F}_{2^p} \times \mathbb{F}_{2^p}} nl(D_{(a,b)}h) \\ & \geq (2^p - 1)(2^{2p-1} - 2^{p+e-1}) + (2^p - 1)(2^{2p-1} - 2^{\frac{3p+e-2}{2}} - 2^{\frac{p+e+2}{2}}) \\ & + (2^p - 1)(2^p - 1)(2^{2p-1} - 2^{p+e-1} - 2^{\frac{p+e+2}{2}}) \\ & = 2^{4p-1} - 2^{3p+e-1} - 2^{2p-1}(1 - 2^e) - 5(2^{\frac{5p+e-2}{2}} - 2^{\frac{3p+e-2}{2}}) \end{aligned}$$



$$nl_2(h) \geq 2^{n-1} - \frac{1}{2} \sqrt{2^{2n} - 2 \sum_{(a,b) \in \mathbb{F}_2^p \times \mathbb{F}_2^p} nl(D_{(a,b)}h)}.$$

Comparisons

$n = 2p$	6	10	12
i	1, 2	1, 2, 3, 4	2, 4
$e = \gcd(i, p)$	1	1	2
Lower bounds in (Gangopadhyay et al.)	15	378	1524
Hamming distances in (Fourquet et al.)	18	400	1760
Lower bounds of D_0 type considered	10	351	1466

Another class of functions

- ▶ Let $h(x, y) = \text{tr}_1^p(x(y^{2^{m+1}+1} + y^3 + y)) + \prod_{i=1}^p(x_i + 1)$, where $n = 2p$, $x, y \in \mathbb{F}_2^p$, m is integer such that $p = 2m + 1$, then

$$nl_2(h) \geq 2^{2p-1} - \frac{1}{2} \sqrt{2^{3p+2} - 3 \cdot 2^{2p} + 5 \cdot (2^{\frac{5p+3}{2}} - 2^{\frac{3p+3}{2}})}.$$

- ▶ S. Sarkar and S. Gangopadhyay, On the Second Order Nonlinearity of a Cubic Maiorana-McFarland Bent Function, Finite Fields and their Applications, Fq 9, Dublin, Ireland, July 13-17, 2009.

Conclusions

- ▶ We identify a class of bent functions, with maximum algebraic degree, having good second order nonlinearity.
- ▶ Finding out bounds of the nonlinearity profile of these functions is an open question.

Acknowledgements

- ▶ The authors thank Projet SECRET, INRIA - Rocquencourt for travel support to present this paper.
- ▶ Brajesh Kumar Singh's research is supported by Council for Scientific and Industrial Research India.

THANK YOU
QUESTIONS PLEASE